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An Efficient Fault Tolerant Job Scheduling and Load Balancing in Grid Computing

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Abstract:

In Grid environment, the two main techniques that are most suitable to cope with the dynamic nature of the grid are load balancing and job replication. In Load balancing algorithm juxtaposes is considered the strong point is neighbour based and cluster based methods. We propose a genetic algorithm for job scheduling to address the heterogeneity of fault-tolerance mechanisms problem in a computational grid. We assume that the system supports four kinds fault-tolerance mechanisms, including the job retry, the job migration without check pointing, the job migration with check pointing, and the job replication mechanisms. Because each fault-tolerance mechanism has different requirements for the gene encoding, we also propose a new chromosome encoding approach to integrate the four kinds of mechanisms in a chromosome. The risky nature of the grid environment is also taken into account in the algorithm. The risk relationship between jobs and nodes are defined by the security demand and the trust level.

Key words: Computational grid, Fault tolerance, Job scheduling, Genetic algorithm

1. Introduction

Grid computing has emerged as the next-generation parallel and distributed computing methodology that aggregates dispersed heterogeneous resources for solving various kinds of large-scale applications in science, engineering and commerce. In large-scale grid environments, the underlying network connecting them is heterogeneous and bandwidth across resources varies from link to link. Not limited to the grid, in many of today's disturbed computing environment (DCE), the computers are linked by a delay and the bandwidth limited communication medium that inherently inflicts tangible delays. Grid systems are classified into two categories: Compute and data grids. In computing grids, the main resources that are being managed by the resources management system is computed cycles, while in data grids the focus is to manage data distributed over geographical locations. The type of grid system, it is deployed in affects the architecture and the services provided by the resources management system. In this paper, we consider fault-tolerant scheduling using genetic algorithm and balancing, application load for a computational grid by taking into account grid architecture.

2. Related Work

Although load balancing, job scheduling, and fault tolerance are active areas of research in a grid environment, these areas have largely been and continue to be developed independent of one another each focusing on different aspects of computing. With respect to load balancing Jasma Balasangameshwara and Nedunchezian Raju [1] proposed That the load balancing with four policies and proved the load adjustment policy is considered as the best policy. Malarvizhi Nandagopal and Rhymend V Uthariaraj [2] proposed load-balancing involves assigning job to a resource proportional to its performance, thereby minimizing the response time of a job. However, there are wide varieties of issues that need to be considered for a heterogenous grid environment. Manimaran and Sivaram Murthy [3] proposed real time-system in which the correctness of the system depends not only on the logical result of computation, but also on the time at which the results are produced. In dynamic scheduling, when a new set of tasks arrive at the system. Qin Zheng, Bharadwaj Veeravalli, Chen-Khong Tham [4] proposed that the need of efficient algorithm that can schedule both independent and dependent jobs on the computational grids so that jobs can run successfully even when processor failure occurs. Besides the need to deal with processor failure, grid applications. Xiaomin Zhu, Xiao Qin,

Meikang Qiu [5] proposed heterogenous clusters, consisting of a diverse computers connected, by high speed networks, have been adopted as an efficient high performance computing platforms for computing-intensive.

3. Load Balancing Policies

In load balancing each resources c_i maintains p number of neighboring resources $LNSet_i$, and cluster resources $LPSet_i$ which the grid scheduler will use to select a neighboring resources for offloading jobs. Neighbors for each resource are formed in terms of transfer delay. For c_i, c_k is considered as its neighbor resources as long as he transfer delay between c_i, c_k is within α times of the transfer.

3.1. Resource Efficiency Estimation Policy

Grid Scheduler maintains the submitted job list, namely

$$S(c_i) = \{-1 \text{ job is finished}\} \\ \{0 \text{ no job}\} \\ \{+1 \text{ job is unfinished}\} \quad (1)$$

Finally the health of the resources is calculated by adding up all the scores.

Grid Scheduler maintains the efficiency level of neighbor-based and cluster-based resources.

3.2. Performance Benefit Factor

Performance benefit factor associated with a job it is based on communication cost taken by the resources.

$$PB_x = cc(j_x, c_i, c_i, t) - cc(j_x, c_i, c_k, t) \quad (3)$$

Where $cc(j_x, c_i, c_k, t)$ is estimated communication cost at c_k

Where $cc(j_x, c_i, c_i, t)$ is estimated communication cost at c_i .

3.3. Enhanced MIF Policy

Each resources c_i maintains the state information about other resources. The information is used to estimate the load and efficiency of other resources. It makes the site process for each resource. Mutual Information Feedback policy is that the rate of load dissemination is directly proportional to the job arrival rate.

3.4. Instantaneous Distribution Policy

When a new job arrives at a resource c_i , the policy decides whether the job has to be executed locally or sent to other resources in $LPSet_i$ $LNSet_i$. The decision depends on whether it can gain performance benefit if it is distributed to a resource in $LPSet_i$ or $LNSet_i$

```

•  $\forall j_x \in J$  with  $origResource(j_x) = c_i \in C$ 
• For each  $c_k$  in  $LPSet_i \cup LNSet_i + c_i$ 
• Calculate  $cc(j_x, c_i, c_k, t)$ 
• Calculate related  $PB_x$ 
• End for
• Find a resources  $c_k$  that gives maximum  $PB_x$ 
• If  $PB_x > 0$  then
•  $c_k$  download  $j_x$ 
• Update Submitted Job List on  $c_i$ 
• Update state object of  $c_k$ 
• EndIf
• (12) Else if  $PB_x \leq 0$ 
•  $Q(c_i) \leftarrow enqueue(j_x)$ 
• End Elself

```

Figure 1: The Pseudocode of instantaneous distribution policy

3.5. Load Adjustment Policy

The load adjustment policy for a resource c_i tries to continuously reduce load difference among c_i and its neighboring resources in $LNSet_i$ by migrating jobs from heavily loaded resources to lightly loaded resources. The Load adjustment policy is triggered whenever c_i receives load information from its neighboring resources. The load-balancing algorithm will use the most recent load status information and efficiency value to initiate a job. Taking into account that all resources are exchanging their loads in parallel and the dynamic nature of the environment, the network reaches a local optimum quickly. Thus, each resource submits some jobs to one of its neighbor resources which have the minimum load among all. If all its neighbors are busier than the resource itself, no job is submitted from the current resource. If there are any jobs in job queue waiting to be executed, the resource tries to submit them to a lighter neighbor resource

```

• Source load
• Sort LNSeti in ascending order based on load
• while running
• do if Q(ci) size > 0
• then
• lightest load
• second HighestLoad
• Velocity
• Threshold
• while velocity>threshold
• do
• submitJobs(velocity,lightestLoadedResources)
• Update State Object of LightestLoadedResources
• sourceLoad
• velocity
• End while
• End doIf
• End while

```

Figure 2: The Pseudocode of load adjustment Policy

4. Fault Tolerance Stratgies

Fault tolerance scheduling policy fault-tolerant scheduling policy namely Minimum Replication Cost (MinRC) and Completion Time for independent and aperiodic jobs is MinRC handles both types of failures that can happen in a grid environment. If a resource shuts down manually, it sends a notice message to its backup and other neighbor and partner resources before shutting down. If a resource fails suddenly without any notice, then the following strategy is employed. We use job replication strategy in which each replica is independently scheduled in a different resource to achieve fault tolerance. The motivation for job replication is that resources in the grid environment can be highly underutilized; thus, these spare resources can be used to run job replicas. MinRC employs peer-to-peer fail-over strategy such that each resource is a backup of another neighbor or partner resource in the grid system

4.1. Boundary Schedules

Replication cost is defined as the actual percentage of time needed to be scheduled for the backup besides all over-loaded periods with existing backups. The neighbor and partner resources for the primary c_i . In case a tie happens, the boundary schedule which can complete earliest is selected.

```

• IfSchedule Eligibility ( $t_s(j_i)$ ) = True
• Then
• cost  $\leftarrow$  Replication Cost
• EndIf
• If cost is less than  $R^R(j_i)$ 
•  $C_B(j_i) \leftarrow C_i$ 
•  $R^R(j_i) \leftarrow$  cost
•  $t^{B,f}(j_i) \leftarrow t_s(j_i) + t_c(j_i)$ 
• EndIf

```

Figure 3: The Pseudocode for Boundary Schedule

4.2. Algorithm for Scheduling Backups

The boundary schedule which has minimum replication cost is chosen the boundary schedule which can complete earliest is selected. Let t_s and t_c denote start time and completion time of an existing schedule, respectively

```

• For  $c_i = c_j$  do
• Invoke Boundary Schedule() with parameters  $t^{B\#p}(j_i)$  and deadline respectively
• For existing schedule within or overlapping with time Window do
• If the existing schedule is primary or non-overloadable backup schedules then
• Invoke Boundary schedule() with parameter  $t_1-t_2(j_i)$  and  $t_r$  respectively
• EndIf
• Else
• Invoke Boundary schedule() with parameter  $t_1-t_2(j_i)$  and  $t_r$  respectively
• End ElseIf
• EndFor
• EndFor

```

Figure 4: The Pseudocode for Scheduling backups

4.3. Schedule Eligibility

A schedule is eligible if it is within the time window and does not overlap with any primary schedule or nonover-loadable backup schedule. Specifically, there must not exist any primary schedule or nonover-loadable backup sche-dule that starts and/or ends in this schedule or contains this schedule. Let t_{lb} and t_{rb} denote start and completion time of overloadable backup schedule.

```

• Initialize  $t^{NO}(c_i)$  to zero
• Initialize  $t^O(c_i)$  to  $t_s(j_i)$ 
• Overloadable backup schedules are browsed in non-decreasing order of their start time
• If  $t^O(c_i)$  is equal  $t^{B\#f}$ 
• Then break
• EndIf
• If  $t_{rb}$  is greater than  $t^O(c_i)$ 
• Then
• If  $t_{lb}$  is greater than  $t^O(c_i)$ 
• Then
•  $t^{NO}(c_i) = t^{NO}(c_i) + t_{lb} - t^O(c_i)$ 
• Initialize  $t^O(c_i)$  to  $t^{B\#f}$ 
• EndIf
• If  $t^O(c_i)$  is greater than or equal  $t^{B\#f}$ 
• Then
• Initialize  $t^O(c_i)$  is equal  $t^{B\#f}$ 
• Break
• EndIf
• EndIf

```

Figure 5: The Pseudocode for Schedule Eligibility

5. Proposed Work

GA-based job scheduling strategy for a large-scale computational grid to address the problem that all existing strategies assume that all computational nodes are protected by the same fault-tolerance mechanism. In fact, it is very common that different autonomous domains usually have different administrative policies and different computational sites have different hardware and software fault-tolerance supports. Therefore, we considered the computational grid in which each computational site supports one or two of four kinds of fault-tolerance mechanisms, including the job retry, the job migration without checkpointing, the job migration with checkpointing, and the job replication mechanisms. The scheduler will decide which kinds of fault-tolerance mechanisms will be applied to each individual job for more reliable computation and shorter makespan. To encode mixed kinds of fault-tolerance mechanisms into a single chromosome, we also proposed a new chromosome encoding approach.

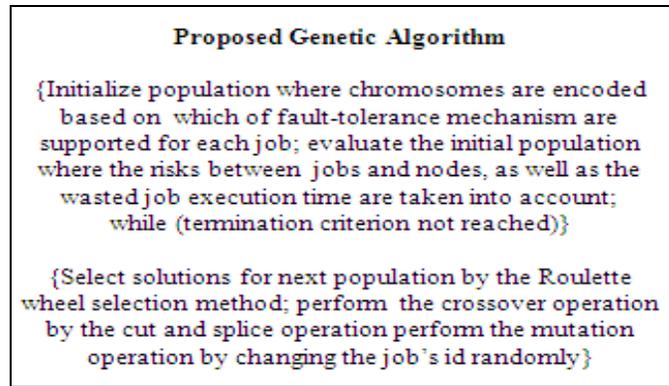


Figure 6: The proposed genetic algorithm for fault-tolerance job Scheduling

- **Job retry (JRT) mechanism**

The JRT mechanism is the simplest fault-tolerance technique, which will re-execute the failed job from the beginning on the same computational node.

- **Job migration without checkpointing (JMG) mechanism**

The JMG mechanism will move the failed job to another computational node and re-execute the job from the beginning on the latter computational node.

- **Job migration with checkpointing (JCP) mechanism.**

The JCP mechanism will record the state of the job periodically at run time. If the job fails, it is moved to another computational node and resumed the execution from the last checkpoint

- **Job Replication (JRP) mechanism**

The JRP mechanism replicates a job to multiple computational nodes such that the job has higher success rate. If one of those replicas has already completed, then all other replicas should stop their execution to save the computing power.

6. Conclusion

In particular, we present an approach for fault tolerance mechanisms as genetic algorithm, validating fault tolerance properties of each mechanism, finding the fitness value for the failed resources and try out different mechanism for failed resources. Therefore we considered the computational grid in which each computational site supports one or two of four kinds of fault-tolerance mechanisms, including the job retry, the job migration without check pointing, the job migration with check pointing, and the job replication mechanism.

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